



## **QFuzz: Quantitative Fuzzing for Side Channels**



### Yannic Noller









# **Detection / Quantification of Side-Channel Vulnerabilities**

#### stringEquals (Original Jetty, v1)

conditional early return causes leakage

#### **Side Channel Vulnerability**

- leackage of secret data
- **software** side-channels
- **observables** (e.g., execution time)

#### **Detection vs Quantification**

Is there a vulnerability? ↔ How much information can be leaked?

# State of the Art

Challenges

How to go beyond non-

interference?

How to avoid **expensive** 

symbolic execution?

How to **scale** to larger programs?

How to provide **guarantees** for

vulnerability?

Blazer

**Decomposition Instead of Self-Composition for** *k*-Safety

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### Precise Detection of Side-Channel Vulnerabilities usingQuantitative Cartesian Hoare Logic

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#### DIFFUZZ: Differential Fuzzing for Side-Channel Analysis

DifFuzz

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## Multi-run side-channel analysis usingMaxLeakSymbolic Execution and Max-SMT

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0 ≤ <sup>3</sup>

# Quantification

#### Information Leakage: min-entropy [Smith2009]

Assuming that the program P is deterministic and the distribution over secret input  $\Sigma$  is uniform, then the information leakage can be characterized  $log_2 k^*$  ( $\epsilon$ =0).

maximum number

of classes in the

#### **Threat Model**

Attacker can pick an **ideal public input** to compromise the **secret value** or some properties of it in **one try**.

#### **Problem Statement**

Find set of secret values  $\Sigma$  and public value y<sup>\*</sup> that characterize the maximum number of observation classes with the highest distance  $\delta$ .

cost observations How to characterize How to identify such inputs? observation classes? fuzzing driv 6 mutated files that show **Partitioning** C(S1, Y) Algorithm a) #partitions D1 b) minimum check for C(S2, y) distance δ P[s<sub>2</sub>, y] mproved partitioning nartition ) program or coverage coverage C(S3, y)**KDynamic &** D2 Fuzzing mutant selection by input evaluation fe Greedy c(s4, y aximize numl of partition

 $log_2 |\Sigma_{Y=y^*}|$ 

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# **Background: Greybox Fuzzing**



the instrumented program P

# **QFuzz: Workflow**



# **KDynamic vs Greedy Partitioning**



# **KDynamic vs Greedy Partitioning**



# **KDynamic vs Greedy Partitioning**



## **Example** (K=100, $\epsilon$ =1, length=16, count=bytecode-instruction)

K=17 δ=3

<pre>stringEquals (Original Jetty, v1) boolean stringEquals(String s1, String s2) {     if (s1 == s2)         return true;     if (s1 == null    s2 == null            s1.length() != s2.length())         return false;     for (int i = 0; i &lt; s1.length(); ++i)         if (s1.charAt(i) != s2.charAt(i))</pre>	<pre>stringEquals (Current Jetty, v4) boolean stringEquals(String s1, String s2) {     if (s1 == s2) return true;     if (s1 == null    s2 == null)         return false;     boolean result = true;     int 11 = s1.length();     int 12 = s2.length();     for (int i = 0; i &lt; 12; ++i)         result &amp;= s1.charAt(i%11) == s2.charAt(i));     return result &amp;&amp; 11 == 12; }</pre>	K=9 δ =1
<pre>stringEquals (Safe Jetty, v5) boolean stringEquals(String s1, String s2) {     if (s1 == s2) return true;     if (s1 == null    s2 == null)</pre>	<pre>Equals (Unsafe Spring-Security) boolean Equals(String s1, String s2) {     if (s1 == null    s2 == null)         return false;     byte[] s1B = s1.getBytes("UTF-8");     byte[] s2B = s2.getBytes("UTF-8");     int len1 = s1B.length;     int len2 = s2B.length;     if (len1 != len2)         return false;     int result = 0;     for (int i = 0; i &lt; len2; i++)         result  = s1B[i] ^ s2B[i];     return result == 0; }</pre>	K=2 $\delta = 149$ only exist special

K=1

DifFuzz

only leaks existence of

special character



#### **Research Questions**

- **RQ1** Which partitioning algorithm (*KDynamic* or *Greedy*) performs better in terms of correct number of partitions and time for partition computation?
- **RQ2** How does QFuzz compare with state-of-the-art SC detection techniques like Blazer, Themis, and DifFuzz?
- **RQ3** Can QFuzz be used for quantification of SC vulnerabilities in real-world Java applications and how does it compare with MaxLeak?

#### Subjects

- Micro-benchmark
- DARPA STAC
- GitHub projects

#### Tools/Techniques

- Blazer
- Themis
- DifFuzz
  - MaxLeak

Our **open-source** tool **QFuzz** and all experimental subjects are **publicly available**:

https://github.com/yannicnoller/qfuzz

http://doi.org/10.5281/zenodo.4722965





# **RQ2 Detection**

Benchmark	Version	QFuzz		DifFuzz	Time (s)			
Denemiaix		<i>p</i> <sub>max</sub>	$\delta_{max}$	$\delta_{max}$	QFuzz, $\overline{p} > 1$	DifFuzz, $\overline{\delta} > 0$	Blazer	Themis
Array	Safe	1	0	1	-	7.40 (+/- 1.21)	1.60	0.28
Array	Unsafe	2	192	195	5.70 (+/- 0.21)	7.40 (+/- 0.93)	0.16	0.23
LoopAndbranch	Safe	2	4	4,278,268,702	1045.33 (+/- 43.51)	18.60 (+/- 6.40)	0.23	0.33
LoopAndbranch	Unsafe	2	4	4,294,838,782	1078.63 (+/- 61.04)	10.60 (+/- 2.62)	0.65	0.16
Sanity	Safe	1	0	0	-	-	0.63	0.41
Sanity	Unsafe	2	3,537,954,539	4,290,510,883	1414.13 (+/- 102.27)	163 (+/- 40.63)	0.30	0.17
Straightline	Safe	1	0	0	-	-	0.21	0.49
Straightline	Unsafe	2	8	8	7.47 (+/- 0.18)	14.60 (+/- 6.53)	22.20	5.30
unixlogin	Safe	-	-	3	-	510 (+/- 91.18)	0.86	-
unixlogin	Unsafe	2	6,400,000,008	3,200,000,008	1784.47 (+/- 21.27)	464.20 (+/- 64.61)	0.77	-
modPow1	Safe	1	0	0	-	-	1.47	0.61
modPow1	Unsafe	22	117	3,068	4.73 (+/- 0.16)	4.80 (+/- 1.11)	218.54	14.16
modPow2	Safe	1	0	9	-	-	1.62	0.75
modPow2	Unsafe	31	1	5,206	294.70 (+/- 104.66)	23.00 (+/- 3.48)	7813.68	141.36
passwordEq	Safe	1	0	0.00	-	-	2.70	1.10
passwordEq	Unsafe	93	2	127	4.57 (+/- 0.22)	8.60 (+/-2.11)	1.30	0.39
k96	Safe	1	0	0	-	-	0.70	0.61
k96	Unsafe	93	2	3,087,339	4.57 (+/- 0.22)	3.40 (+/- 0.98)	1.29	0.54
gpt14	Safe	12	1	517	5.00 (+/- 0.00)	4.20 (+/- 0.80)	1.43	0.46
gpt14	Unsafe	92	2	12,965,890	5.87 (+/- 0.12)	4.40 (+/- 1.03)	219.30	1.25
login	Safe	1	0	0	-	-	1.77	0.54
login	Unsafe	17	2	62	7.77 (+/- 0.69)	10.00 (+/- 2.92)	1.79	0.70

same vulnerabilities detected

additional information about the strength of leaks and the exploitability of vulnerabilities

**large** values for *K* may slow down **QFuzz**, but eventually, enable the exploration of many partitions

# **RQ3 Quantification**

Modulo Lon		#Partitions	QFuzz ( $\epsilon$ =0, 1h)				MaxL	eak (default)	MaxLeak (No solver)	
	$\overline{p}$		$p_{max}$	Time (s): $p_{max}$	t <sub>min</sub>	#Obs	Time (s)	#Obs	Time (s)	
1717	3	7	7.00 (+/- 0.00)	7	1.00 (+/- 0.00)	1	6	20.892	9	1.047
1717	4	10	10.00 (+/- 0.00)	10	7.43 (+/- 0.45)	5	9	152.332	12	1.370
1717	5	13	13.00 (+/- 0.00)	13	20.40 (+/- 3.87)	6	12	839.788	15	2.916
1717	6	16	16.00 (+/- 0.00)	16	294.60 (+/- 53.17)	22	15	3731.328	18	8.006
1717	7	19	18.37 (+/- 0.25)	19	2484.30 (+/- 451.42)	385		> 4 h	21	19.241
1717	8	22	20.43 (+/- 0.45)	22	3168.07 (+/- 303.47)	508		> 4 h	24	91.821
1717	9	25	22.20 (+/- 0.36)	24	3489.03 (+/- 169.19)	1009	> 4 h		> 8 GB	
1717	10	28	24.40 (+/- 0.49)	27	3548.63 (+/- 57.73)	2929		> 4 h		> 8 GB
834443	3	7	7.00 (+/- 0.00)	7	13.40 (+/- 1.96)	8	6	7.416	9	1.188
834443	4	10	10.00 (+/- 0.00)	10	40.33 (+/- 12.14)	6	9	42.684	12	1.385
834443	5	13	12.93 (+/- 0.09)	13	645.70 (+/- 329.43)	74	12	215.929	15	2.953
834443	6	16	15.40 (+/- 0.20)	16	2711.87 (+/- 433.23)	271	15	936.921	18	7.511
834443	7	19	16.80 (+/- 0.33)	18	3227.60 (+/- 275.29)	952	18	4021.150	21	19.068
834443	8	22	17.93 (+/- 0.54)	22	3556.70 (+/- 83.44)	2301		> 4 h	24	96.360
834443	9	25	20.13 (+/- 0.59)	24	3572.83 (+/- 37.16)	3110	> 4 h		> 8 GB	
834443	10	28	21.83 (+/- 0.46)	24	3504.13 (+/- 121.70)	1845		> 4 h		> 8 GB
1964903306	3	7	6.47 (+/- 0.18)	7	2228.30 (+/- 542.13)	119	6	12.167	9	1.085
1964903306	4	10	8.67 (+/- 0.19)	10	3494.30 (+/- 203.69)	429	9	70.805	12	1.535
1964903306	5	13	10.70 (+/- 0.19)	12	3594.00 (+/- 11.56)	3420	12	2306.261	15	3.391
1964903306	6	16	12.90 (+/- 0.11)	13	1337.90 (+/- 443.89)	206		> 4 h	18	7.506
1964903306	7	19	14.10 (+/- 0.27)	15	2984.67 (+/- 362.05)	503		> 4 h	21	19.486
1964903306	8	22	15.33 (+/- 0.36)	17	3398.37 (+/- 204.45)	1411		> 4 h	24	98.325
1964903306	9	25	16.30 (+/- 0.51)	19	3562.33 (+/- 54.24)	2819		> 4 h		> 8 GB
1964903306	10	28	17.30 (+/- 0.48)	20	3559.67 (+/- 77.72)	2390		> 4 h		> 8 GB

due to its dynamic analysis, **QFuzz** is **more scalable** than **MaxLeak** 

**QFuzz** has precision comparable to **MaxLeak** that uses symbolic execution with model counting

even for complex scenarios **QFuzz** provides reasonable **lower-bound guarantees** 

